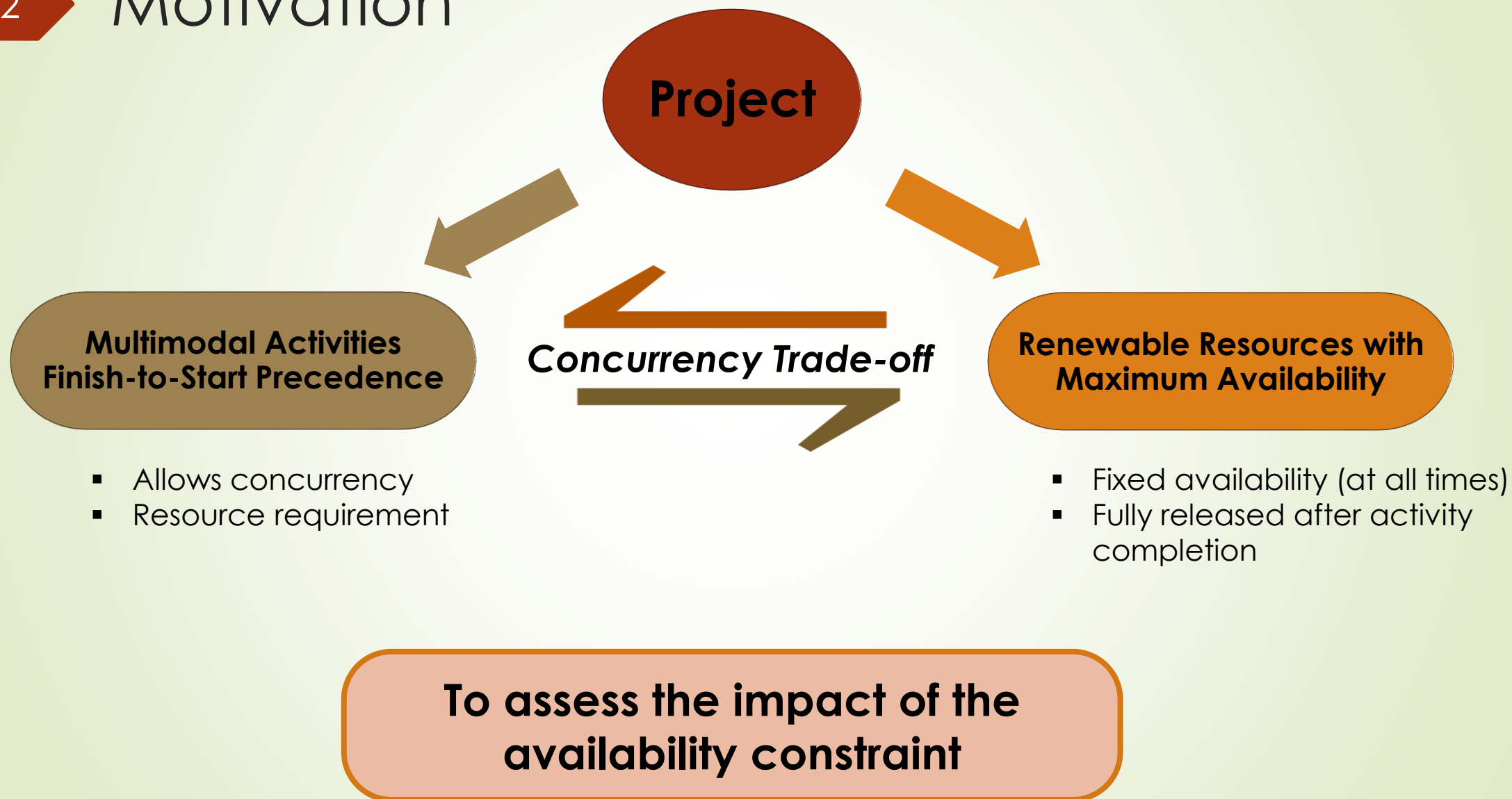


Concurrency Detection on Finish-to-Start Activity Precedence Networks

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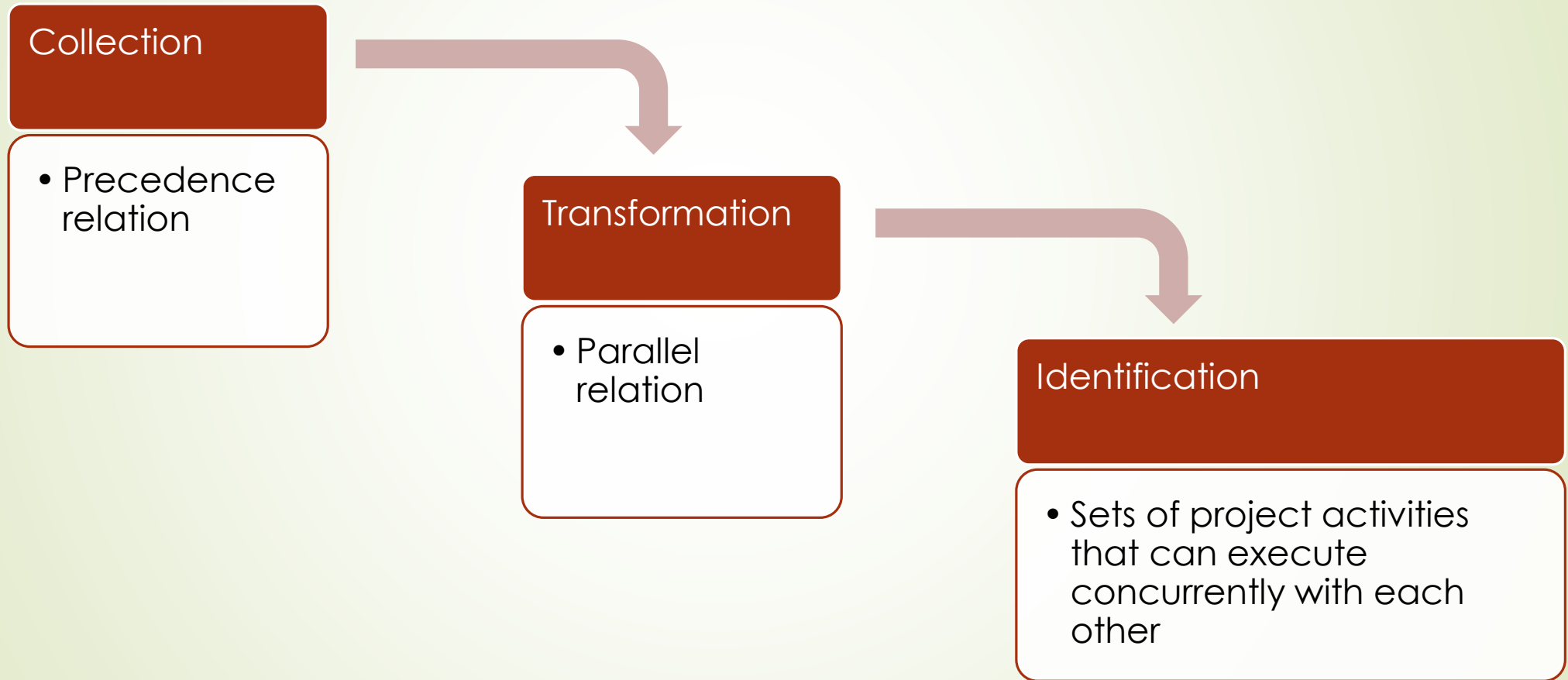
2

Motivation



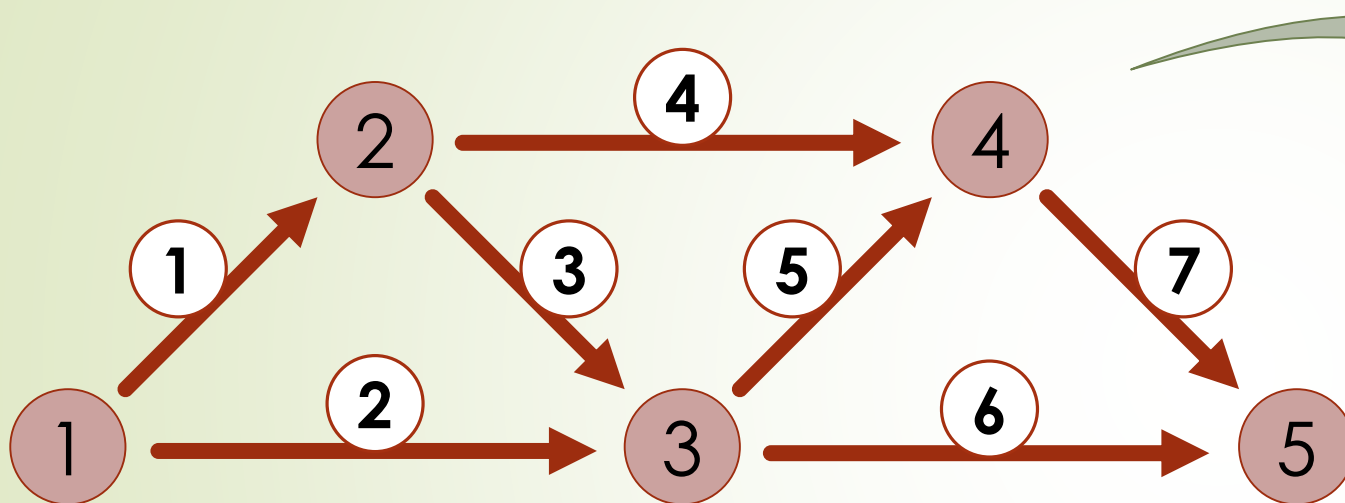
3

Concurrency Detection Method



4

Detection Method: Collection



Finish-to-Start Activity Precedence Relation

usually specified through either AoA or AoN networks

(here, in AoA: 7 arcs – 7 activities)



Intuitively **Transitive**

	1	2	3	4	5	6	7
1							
2							
3							
4							
5							
6							
7							

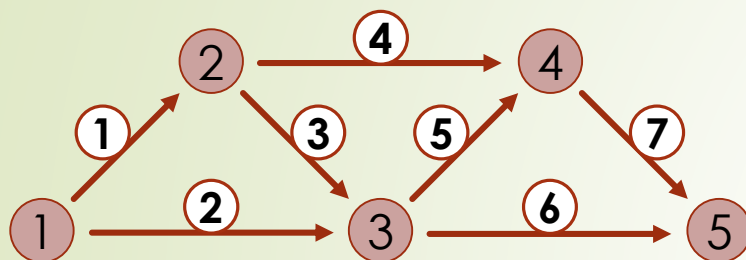
Finish-to-Start Activity Precedence (matrix form)

row x column = reads:

“row must precede column”

5

Detection Method: Transformation



	1	2	3	4	5	6	7
1			■	■	■	■	■
2					■	■	■
3					■	■	■
4							■
5							■
6							
7							

Transitive Precedence Relation

$row \times column = \blacksquare$ means

"row" **must occur before** "column"



If neither activity a or b precedes the other they can execute concurrently

	1	2	3	4	5	6	7
1							
2							
3							
4							
5							
6							
7							

Parallel Relation (matrix form)

$row \times column = \blacksquare$ means

"row" can execute in parallel with "column"

6

Detection Method: Identification

Parallel Relation Matrix

adjacency

Simple Undirected Graph

Set of activities, such:

Set of nodes, such:

Any two different activities may execute in parallel

Any two different nodes are directly connected



There is no activity outside the set which can occur in parallel to any from the set

There is no node outside the set which is directly connected to all in the set



The largest possible set whose activities can execute in parallel

The largest possible set of nodes which are connected to each other



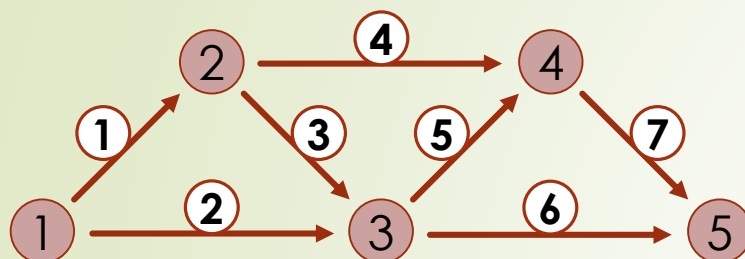
Parallel Set

Maximal Clique

All Parallel Sets \Leftrightarrow All Maximal Cliques

7

Detection Method: Identification



Parallel Sets:

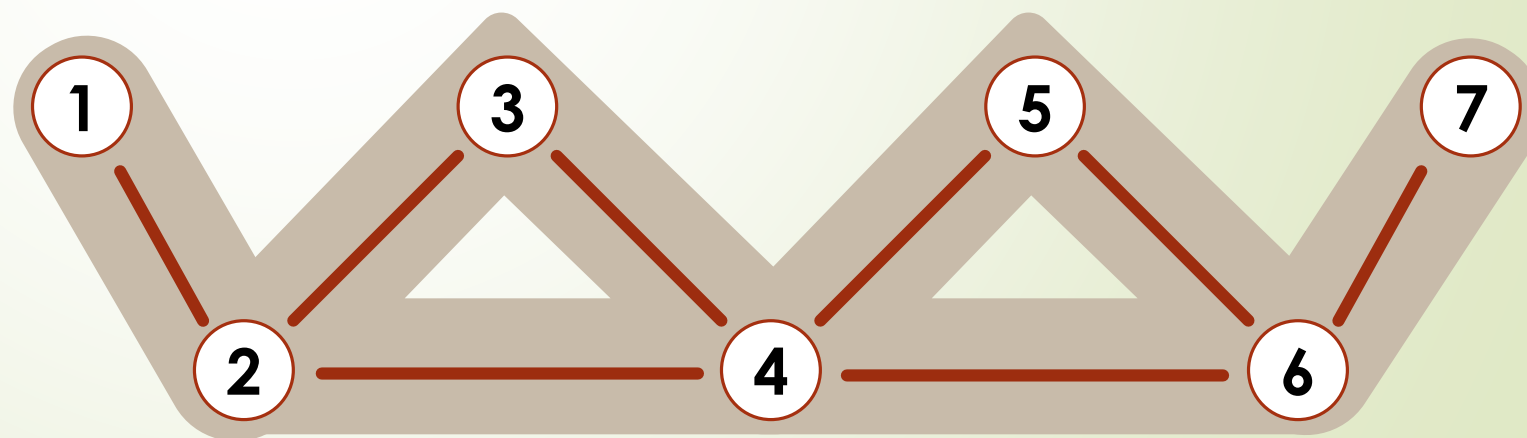
{1, 2}

{4, 5, 6}

{2, 3, 4}

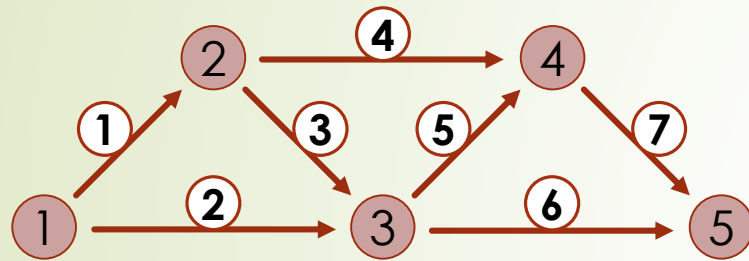
{6, 7}

	1	2	3	4	5	6	7
1							
2							
3							
4							
5							
6							
7							



8

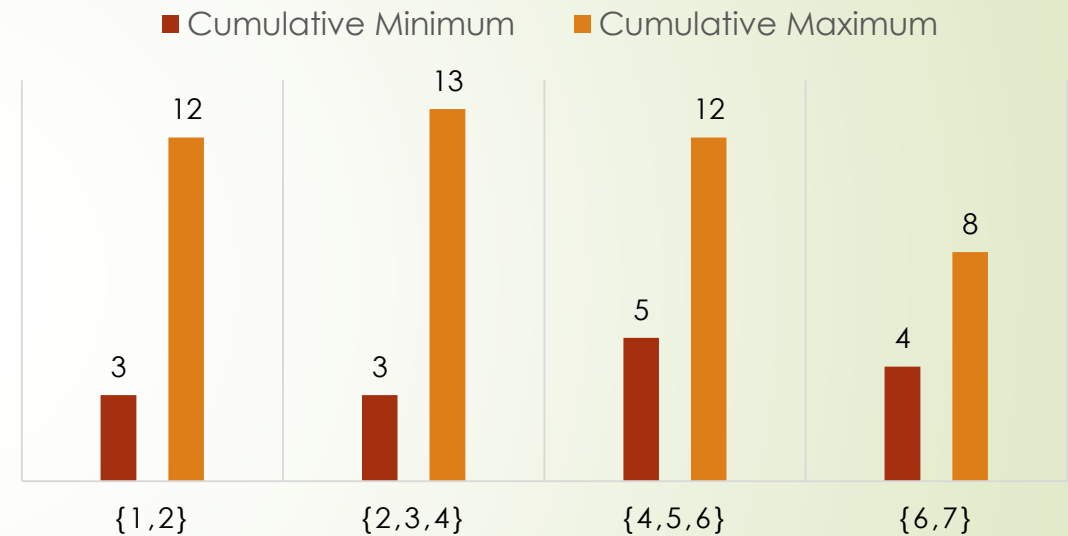
Application: Resource Availability Constraint



Multimodal Activities:

1. Min = 2; Max = 7
2. Min = 1; Max = 5
3. Min = 1; Max = 4
4. Min = 1; Max = 4
5. Min = 2; Max = 3
6. Min = 2; Max = 5
7. Min = 2; Max = 3

PARALLEL SET REQUIREMENTS



- With 13 units or more, no resource constraint is imposed
- With 5 units or more, any set is supported
- With less than 3 units, no concurrency will be possible

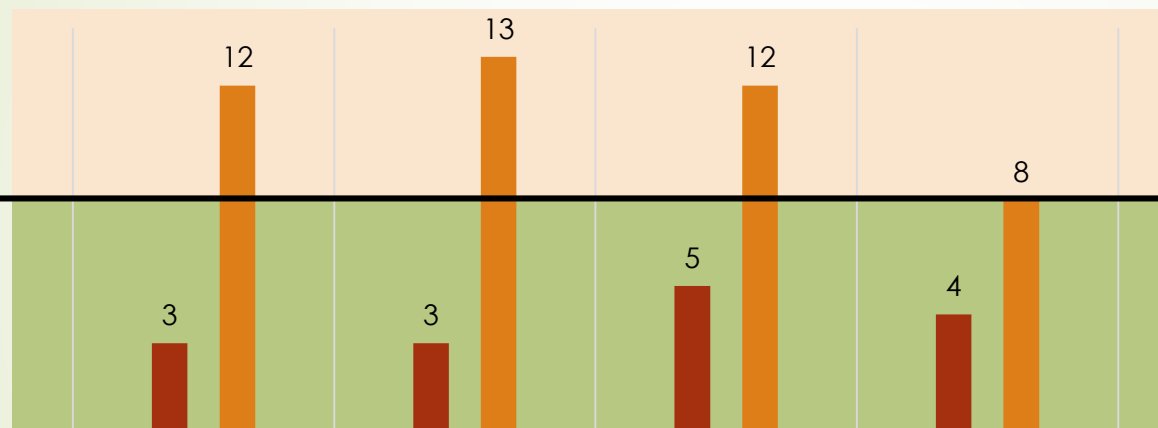
Application: Resource Stress

PARALLEL SET REQUIREMENTS

m Cumulative Minimum
 M Cumulative Maximum
 M

Maximum Availability:
8 units

K



$\{1,2\}$ stress ≈ 0.444
 $\{2,3,4\}$ stress = 0.5
 $\{4,5,6\}$ stress ≈ 0.571
 $\{6,7\}$ stress = 0

Availability **prohibits** concurrency scenarios

Availability **allows any** concurrency scenarios



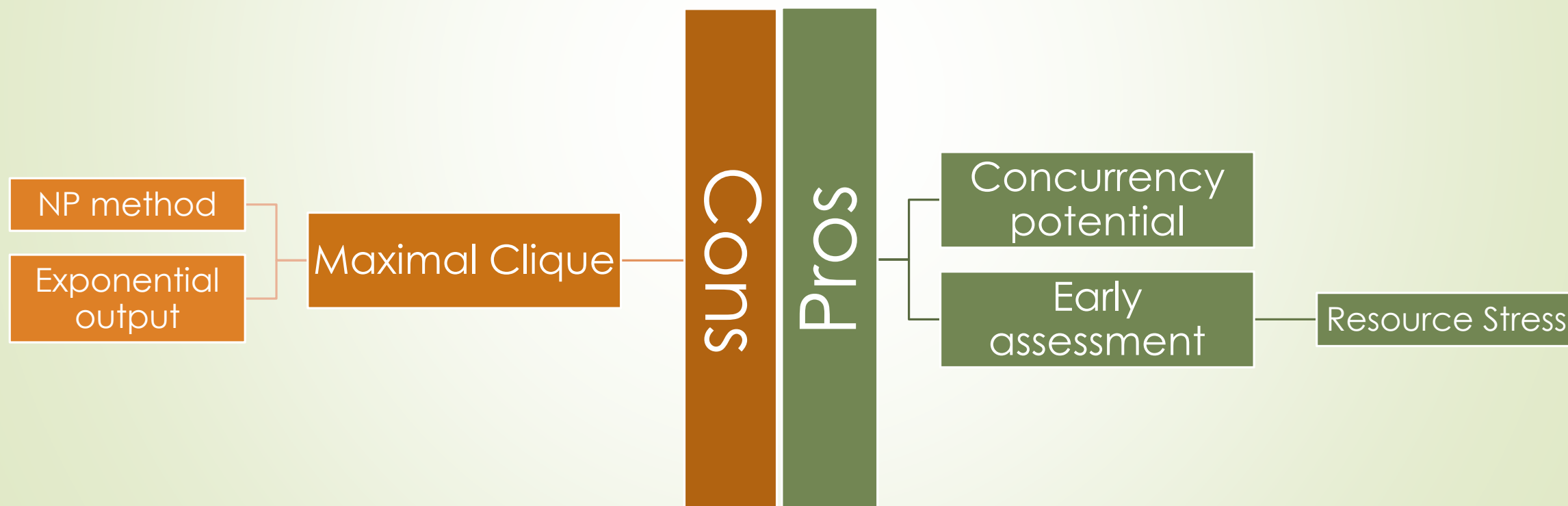
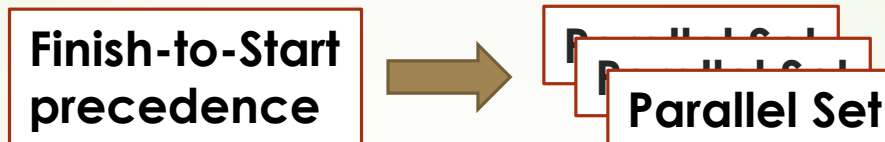
The maximum stress is $\frac{8-8}{8-4} = 0.571$ (57%)

$$\text{stress} = \begin{cases} \frac{M-K}{M-m} & m \leq K \leq M \\ 0 & K < M \\ 1 & K < m \end{cases}$$

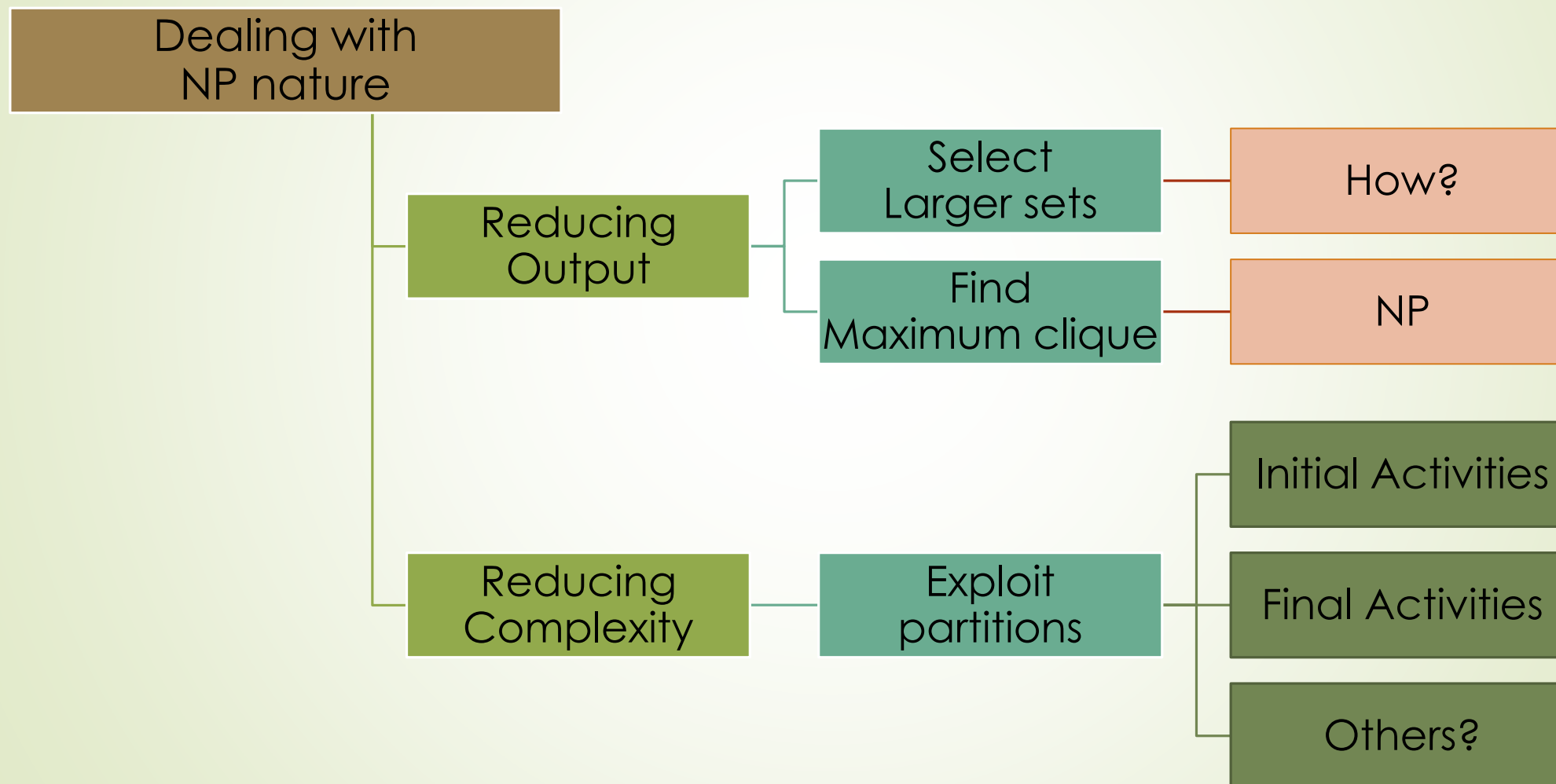
Resource Stress (of a parallel set)

The probability of the cumulative requirement being greater than the maximum availability

Conclusions & Future Research



Conclusions & Future Research



References

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